

Pattern P = sip

i  
ippi  
issippi  
issippi  
issippi  
issippi  
ppi  
ppi  
sippi  
issippi  
issippi  
issippi

11 1  
8 2  
5 3  
2 4  
1 5  
10 6  
9 7  
7 8  
4 9  
6 10  
3 11

$\lfloor \frac{11}{2} \rfloor$  sip > pi

$5 + \lfloor \frac{6}{2} \rfloor$

sip ist Präfix  
von sippi ✓

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Suche durch Intervallhalbierung

S = mississippi

1 mississippi  
2 ississippi  
3 sissippi  
4 sissippi  
5 sissippi  
6 sippi  
7 sippi  
8 ippi  
9 ppi  
10 pi  
11 i

i  
ippi  
issippi  
issippi  
issippi  
mississippi  
pi  
ppi  
sippi  
issippi  
sippi  
ssissippi

Suffix

Array

# Burrows-Wheeler transform

final char ( <i>L</i> )	sorted rotations
a o o o o a a i i o a e i e e i i i o o	<p>n to decompress. It achieves compression</p> <p>n to perform only comparisons to a depth</p> <p>n transformation} This section describes</p> <p>n transformation} We use the example and</p> <p>n treats the right-hand side as the most</p> <p>a n tree for each 16 kbyte input block, enc</p> <p>a n tree in the output stream, then encodes</p> <p>i n turn, set <math>L[i]</math> to be the</p> <p>i n turn, set <math>R[i]</math> to the</p> <p>o n unusual data. Like the algorithm of Man</p> <p>a n use a single set of probabilities table</p> <p>e n using the positions of the suffixes in</p> <p>i n value at a given point in the vector <math>R</math></p> <p>e n we present modifications that improve t</p> <p>e n when the block size is quite large. Ho</p> <p>i n which codes that have not been seen in</p> <p>i n with <math>sch</math> appear in the {\em same order</p> <p>i n with <math>sch</math>. In our exam</p> <p>o n with Huffman or arithmetic coding. Bri</p> <p>o n with figures given by Bell~\cite{bell}.</p>

Auszug aus der BWT Matrix für die Latex Sourcen des Papers  
 In der L-Spalte (eigentlich ganz rechts) finden sich nur Vokale